We also have a convergence results for the approximation of the eigenvalue  $\lambda_1$ .

## **Corollary**

Let  $A \in \mathbb{C}^{n \times n}$  be a diagonalizable matrix. Assume  $|\lambda_1| > |\lambda_2|$  and  $v_0 = \sum_{i=1}^n \alpha_i u_i$ , with  $\alpha_1 \neq 0$ . Then it holds

$$
|\mu_k - \lambda_1| = O\left(\left|\frac{\lambda_2}{\lambda_1}\right|^k\right)
$$
, for  $k \to +\infty$ .

For symmetric real matrices, we have a better convergence results:

### **Corollary**

Let  $A \in \mathbb{R}^{n \times n}$  be a symmetric matrix. Assume  $|\lambda_1| > |\lambda_2|$  and  $v_0 = \sum_{i=1}^n \alpha_i u_i$ , with  $\alpha_1 \neq 0$ . Then it holds

$$
|\mu_k - \lambda_1| = O\left(\left|\frac{\lambda_2}{\lambda_1}\right|^{2k}\right)
$$
, for  $k \to +\infty$ .

## Some observations

One of the hypothesis of the previous results is  $\alpha_1 \neq 0$ , where  $\alpha_i$  are defined such that  $\mathsf{v}_0 = \sum_{i=1}^n \alpha_i \mathsf{u}_i$ . Clearly,  $\mathsf{u}_1, \ldots, \mathsf{u}_n$  are unknown and we cannot check if  $v_0$  satisfies this hypothesis.

Practically this is not a real obstacle. If unfortunately we choose  $v_0$  s.t  $\alpha_1 = 0$ , there are two possible cases:

- in exact arithmetic, we get  $\lim_{k\to+\infty} \tilde{v}_k = u_2$  and  $\lim_{k\to+\infty} \mu_k = \lambda_2$ , as long as  $|\lambda_2| > |\lambda_3|$  and  $\alpha_2 \neq 0$ .
- in *finite arithmetic*, during the iterations of the Power Method, round-off errors can cause the appearance of a non-zero component in the direction of  $u_1$  in a certain  $v_k$ . If this happens, the method immediately starts to converge towards the dominant eigenvalue  $\lambda_1$ and its corresponding eigenvector  $u_1$ .

# Stopping criterion

A simple stopping criterion for the power method is based on the residual:

$$
\text{Stop when } \|Av_k - \mu_k v_k\| \leq \text{tol}
$$

## How can we compute other eigenvalues and eigenvectors?

Let  $\mu \in \mathbb{C}$  a user-specified parameter that is not an eigenvalue of A, we want to approximate the closest eigenvalue of A to  $\mu$ , i.e.

$$
\lambda_J = \underset{i}{\text{argmin}} \left| \mu - \lambda_i \right|
$$

### Inverse Power method

Input:  $A \in \mathbb{C}^{n \times n}$ ,  $v_0 \in \mathbb{C}^n$  with  $||v_0|| = 1$ , MAXITER  $\in \mathbb{N}$ , tol  $\in \mathbb{R}^+$ . for  $k = 1, 2, \ldots$ , MAXITER  $w = (A - \mu I)^{-1} v_{k-1}$  (equivalently, solve  $(A - \mu I) w = v_{k-1}$ )  $v_k = w / ||w||$  $\mu_k = (v_k)^H A v_k$ (Rayleigh quotient with  $A$ ) Check the Stopping criterion end Output:  $\mu_k$  and  $v_k$ .

#### Beware

Since  $\mu$  is not an eigenvalue of A, the matrix  $A - \mu I$  is non singular.

Since  $Au_i = \lambda_i u_i$ , then  $(A - \mu I)u_i = (\lambda_i - \mu)u_i$ , and then  $\frac{1}{\lambda_i-\mu}u_i=(A-\mu I)^{-1}u_i$ . Let  $\lambda_J$  be the eigenvalue of  $A$  closest to  $\mu$ , the largest (in module) eigenvalue of  $(A - \mu I)^{-1}$  is then  $\frac{1}{\lambda_J - \mu}$ , and the relative eigenvector is  $u_1$ . The inverse power method is just a power method applied to  $(A - \mu I)^{-1}$ , and the previous results apply:  $\widetilde{v}_k$  converges to  $u_j$ .<br>Since the Payleigh quotient  $u_k$  is computed with A instead of  $(A - \mu I)^{-1}$ . Since the Rayleigh quotient  $\mu_k$  is computed with  $A$  instead of  $\left(A-\mu I\right)^{-1}$ , it converges to  $\lambda_1$ .

#### Theorem

Assume  $|\mu - \lambda_J| < |\mu - \lambda_i|$   $\forall i = 1, \ldots, n$ ,  $i \neq J$  and  $v_0 = \sum_{i=1}^n \alpha_i u_i$ , with  $\alpha$ ,  $\neq$  0. Then

$$
\lim_{k \to +\infty} \mu_k = \lambda_j
$$

and

$$
\lim_{k \to +\infty} \|\widetilde{v}_k - u_J\|_2 = 0, \qquad \text{where } \widetilde{v}_k = \frac{\|A^k v_0\|}{\alpha_1 \lambda_1^k} v_k.
$$

Note that if  $\mu = 0$ , the method approximates the eigenvalue of A that is smallest in module.