# Model Order Reduction in support of the Virtual Element Method

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# Outline

- The Virtual Element Method
- Parametric PDEs and Reduced Basis Method
- $\bullet~\mathsf{RB}$  in support of VEM
- Examples

# Model problem

### Poisson equation

Find 
$$u \in V = H_0^1(\Omega)$$
 such that

$$a(u,v) = \int_{\Omega} \nabla u \cdot \nabla v \, d\mathsf{x} = \int_{\Omega} fv \, d\mathsf{x} \quad \forall v \in V$$
 (1)

- $\Omega \subset R^2$  is a polygonal domain
- $\Omega$  is decomposed with a mesh  $\mathcal{T}_h$  made up of polygonal elements K



# VEM space, degree 1

- r = 1, polynomial degree (accuracy)
- As in Finite Elements, we define the space in each element

Local VEM space

$$V_1^{\mathcal{K}} = \{ v \in H^1(\mathcal{K}) : v_{|e} \in P_1(e) \ \forall e \in \partial \mathcal{K} \text{ and } \Delta v_{|\mathcal{K}} = 0 \}$$

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- Remark:  $P_1(K) \subset V_1^K$
- DOFs: the values at the vertices of K

We glue them by continuity

### Global VEM space

$$V_h = \{ v \in H^1_0(\Omega) : v_{|K} \in V_1^K \ \forall K \in \mathcal{T}_h \}$$

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### Discretization

**Galerkin method** Find  $u_h \in V_h$  such that

$$a(u_h,v_h) = \int_{\Omega} f v_h \quad \forall v_h \in V_h$$

In particular,

$$a(u_h, v_h) = \int_{\Omega} \nabla u_h \cdot \nabla v_h \, d\mathsf{x} = \sum_{K \in \mathcal{T}_h} \int_{K} \nabla u_h \cdot \nabla v_h \, d\mathsf{x} = \sum_{K \in \mathcal{T}_h} a^K(u_h, v_h)$$

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### Warning!

We cannot assemble  $a^{K}$  on each element directly using the basis functions of  $V_{1}^{K}$  (like in FEM) because they are themselves solutions of PDEs in K

## A computable bilinear form

- Decomposition of  $V_1^K = P_1(K) \oplus V_\perp^K$
- $V_{\perp}^{K} \subset V_{1}^{K}$   $H^{1}-$ orthogonal to  $P_{1}(K)$
- $u_h = p + p^{\perp}$  and  $v_h = q + q^{\perp}$

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- $u_h = p + p^{\perp}$  and  $v_h = q + q^{\perp}$
- p, q computable,  $p^{\perp}, q^{\perp}$  not computable

$$a^{K}(u_{h},v_{h})=a^{K}(p,q)+\underline{a^{K}(p,q^{\perp})}+\underline{a^{K}(p^{\perp},q)}+a^{K}(p^{\perp},q^{\perp})$$

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**Idea:** stabilization term  $S^{\kappa}(p^{\perp},q^{\perp})$  symmetric and bilinear s. t.

$$c_{\star}a^{K}(p^{\perp},p^{\perp}) \leq S^{K}(p^{\perp},p^{\perp}) \leq C^{\star}a^{K}(p^{\perp},p^{\perp})$$

Scaled diagonal:  $S^{K}(p^{\perp},q^{\perp}) = \sum_{i} dof_{i}(p^{\perp}) dof_{i}(q^{\perp}) |\Pi_{k}^{\nabla} e_{i}|_{1,K}$ 

# A look at the Reduced Basis method

# Parametric PDEs and Reduced Basis Method

### Model problem

Find  $u(\mu) \in V$  such that

$$\mathsf{a}(u(\mu), \mathbf{v}; \mu) = f(\mathbf{v}; \mu) \quad \forall \mathbf{v} \in V$$

(2)

where  $\mu$  is a parameter.

• Galerkin. We can compute an approximation

$$u_{\mathcal{N}} \in V_{\mathcal{N}} = span\{\varphi_1, \dots, \varphi_{\mathcal{N}}\} \subset V$$

• if we need to solve this for many values of  $\mu$ , this will be extremely expensive

### Reduced Basis solution

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- $\bullet\,$  We want to introduce a new space  ${\it W}_{\it M} \subset {\it V}_{\it N}$

$$W_M = span\{\xi_1, \ldots, \xi_M\}$$
 with  $M \ll \mathcal{N}$ 

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$$u_M(\mu) = \sum_{i=1}^M u_i \xi_i$$

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 $a(u_M, \xi_j; \mu) = \sum_{i=1}^M a(\xi_i, \xi_j; \mu) u_i(\mu) = f(\xi_j; \mu) \text{ for } i \le j \le M$  (3)

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### How do we construct $W_M$ ?

It is generated by combinations of snapshots  $u(\mu_i)$  for a small set of parameters  $\mu_1, \ldots, \mu_{N_s}$ .

# Computation of $a(\xi_i, \xi_j; \mu)$

### We assume that

• 
$$a(u,v;\mu) = \sum_{q=1}^{Q_a} \theta_q^a(\mu) a^q(u,v)$$

• 
$$f(\mathbf{v};\mu) = \sum_{q=1}^{Q_f} \theta_q^f(\mu) f^q(\mathbf{v})$$

Hence, (3) becomes

$$\sum_{i=1}^{M} \left[ \theta_{q}^{a}(\mu) a^{q}(\xi_{i},\xi_{j}) \right] u_{i}(\mu) = \sum_{q=1}^{Q_{f}} \theta_{q}^{f}(\mu) f^{q}(\xi_{j})$$
(4)

### Offline-online strategy

#### **Offline - Sample**

Build sample  $S = \{\mu_1, \ldots, \mu_{N_s}\}$ 

#### Comment

*S* set of the parameters to build the reduced basis.

- Random
- Equidistributed/log equidistributed
- From error estimator

## Offline-online strategy

Offline - Sample

Build sample  $S = \{\mu_1, \ldots, \mu_{N_s}\}$ 

### Offline - Build basis

 $\forall \mu_i \in S$ 

• 
$$A_{\ell,k}^{\mathcal{N}} = a(\varphi_{\ell}, \varphi_k; \mu_i)$$

• 
$$F_k^{\mathcal{N}} = f(\varphi_k; \mu_i)$$

 Choose the r.b. functions (POD)

#### Comment

Compute the snapshots

- Compute  $u^{\mathcal{N}}(\mu_i) \quad \forall i$
- $\bullet~$  Cost depends on  ${\cal N}$
- $\xi_i$  stored

Only ONCE!

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#### **Offline - Precomputations**

• 
$$A^{M,q}_{i,j} = a^q(\xi_i,\xi_j)$$

• 
$$F_j^{M,q} = f^q(\xi_j)$$

#### Comment

 $\xi_i, \xi_j$  known functions in  $V_N$ Building block for affine decomposition  $a(\cdot, \cdot; \mu)$  and  $f(\cdot, \mu)$ 

• 
$$A_{i,j}^{M,q} = a^q(\xi_i,\xi_j)$$

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Precomputed and stored ONCE!

## Offline-online strategy

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Build sample  $S = \{\mu_1, \ldots, \mu_{N_s}\}$ 

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 $\forall \mu_i \in S$ 

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$$A_{\ell,k}^{\mathcal{N}} = a(\varphi_{\ell}, \varphi_k; \mu_i)$$

• 
$$F_k^{\mathcal{N}} = f(\varphi_k; \mu_i)$$

• Solve 
$$A^{\mathcal{N}}u^{\mathcal{N}}(\mu_i) = F^{\mathcal{N}}$$

 Choose the r.b. functions (POD)

#### **Offline - Precomputations**

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$$A^{M,q}_{i,j} = a^q(\xi_i,\xi_j)$$

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$$F_j^{M,q} = f^q(\xi_j)$$

### **Online:**

For each new parameter  $\mu$ :

• 
$$A_{i,j}^M = \sum_{q=1}^{Q_a} \theta_a^q(\mu) A_{i,j}^q$$

• 
$$F_j^M = \sum_{q=1}^{Q_f} \theta_f^q(\mu) F_j^q$$

• Solve 
$$A^M x = F^M$$

• 
$$u_M(\mu) = \sum_{i=1}^M x_i \xi_i$$

# Proper Orthogonal Decomposition

Build the correlation matrix of all the snapshots of the sample

$$C = rac{1}{N_s} U^T U$$
 where  $U = \begin{bmatrix} u^{\mathcal{N}}(\mu_1) & | & \dots & | & u^{\mathcal{N}}(\mu_{N_s}) \end{bmatrix}$ 

- **2** Solve the eigenvalues problem  $Cz_n = \lambda_n z_n$
- Ochoose the eigenvectors corresponding to the first M greatest eigenvalues
- The POD basis can be computed as

$$\xi_i(x) = \frac{1}{\sqrt{N_s}} \sum_{m=1}^{N_s} (z_n)_m u^{\mathcal{N}}(\mu_m)(x) \quad i = 1, \dots, M$$

# Back to VEM

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## Order 1 VEM

- Consider a polygon K
- $v_1, \ldots, v_N$  vertices of K

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# Order 1 VEM

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# Basis of $V_1^K$ For n = 1, ..., N $-\Delta e^n = 0$ in K $e^n$ p.w lin. on $\partial K$ (5) $e^n(v_m) = \delta_{n,m}$

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# Geometric parametrization

- $\hat{K}$  reference (regular) polygon
- $\mathcal{B}_{\mathcal{K}}:\mathcal{K} 
  ightarrow \hat{\mathcal{K}}$  piecewise affine transformation such that

$${\mathcal B}_{\mathcal K}(v_n) = \hat v_n$$
 and  ${\mathcal B}_{\mathcal K}(x_{\mathcal K}) = \hat x_{\mathcal K}$ 

for  $x_K \in K$  and  $\hat{x}_K \in \hat{K}$ .

• In particular,  $B_K$  is a piecewise constant matrix



# Affine decomposition

• We can partition K and  $\hat{K}$  in as many triangles as there are edges

$$\hat{K} = \cup_{n=1}^{N} \hat{T}_n$$
 and  $K = \cup_{n=1}^{N} T_n$ 

where  $\hat{T}_n = \mathcal{B}_K T_n$ 

•  $\mathcal{B}_K$  is affine on  $T_n$ , we have

$$a(u,v;K) = \sum_{n=1}^{N} \int_{\hat{T}_n} B_K B_K^T \nabla \hat{u} \cdot \nabla \hat{v}$$

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$$\mathsf{a}(u,v;K) = \sum_{n=1}^{N} \int_{\hat{T}_n} B_K B_K^T \nabla \hat{u} \cdot \nabla \hat{v}$$

• 
$$B_K B_{K|T_n}^T = \sum_{\nu=1}^3 a_{\nu}^n A^{\nu} = a_1^n \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} + a_2^n \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} + a_3^n \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$
  
• Hence,  $q = (n, \nu) \Rightarrow a_{\nu}^q (u, \nu) := \int_{\hat{T}_n} A^{\nu} \nabla u \cdot \nabla v$ 

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# Offline phase

- N is fixed
- Generate a set of trial polygons  $K^{\ell}$
- Compute the affine mapping  $\mathcal{B}_\ell: \mathcal{K}^\ell o \hat{\mathcal{K}}$
- Compute  $e^1_\ell, \ldots, e^N_\ell$  by solving their equations in  $\mathcal{K}^\ell$  (FEM)
- Map on  $\hat{K}$  the VEM basis just computed  $\rightsquigarrow \hat{e}^1_\ell, \dots, \hat{e}^N_\ell$
- Compute and store  $a^q_{
  u}(\hat{e}^n_\ell,\hat{e}^{n'}_\ell)$

# Online phase

After building the reduced basis  $\{\hat{\xi}^n_\ell\in H^1(\hat{K}):\ell=1,\ldots,M\}$  using the POD

- Generate a set of test polygons
- $\forall K, \mathcal{B}: K \to \hat{K} \text{ s.t. } BB_{|T_i}^T = \sum_{\nu=1}^3 a_i^{\nu} A^{\nu}$
- We look for  $\hat{e}^n = \sum_{\ell=1}^M x_\ell^n \hat{\xi}_\ell^n$  s.t.

$$-\nabla \cdot (BB^{T})\nabla \hat{e}^{n} = 0 \text{ in } \hat{K}$$
$$\hat{e}^{n} \text{ p.w lin. on } \partial \hat{K}$$
$$\hat{e}^{n}(\hat{v}_{m}) = \delta_{n,m}$$

 ê<sup>n</sup> are the VEM basis for K mapped on K
 , hence we go back
 to obtain e<sup>n</sup> for n = 1,..., N.

# Different uses of our basis

- Build the stabilization matrix
- Post-processing of VEM solutions and reconstruction in subdomains
- Evaluate the error with respect to the true solution (academic purpose)

### Reduced Basis generation for VEM stabilization

- In order to perform a numerical test on a VEM solution, we generated Reduced Basis (with several choices of *M*) on sets of convex random polygons with N = 4, 5, ..., 14
- For each *N*, we generated 300 trial polygons and 500 test polygons
- We studied the ratio  $C^{\star}/c_{\star}$ , where

$$c_{\star}a^{K}(p^{\perp},p^{\perp}) \leq S^{K}_{RB}(p^{\perp},p^{\perp}) \leq C^{\star}a^{K}(p^{\perp},p^{\perp})$$

# N = 6 - some polygons



## N = 11 - some polygons



# Ratio $C^*/c_*$



# RB for post-processing

- We want to solve Poisson for  $u(x, y) = \frac{\sin(4x\pi)\sin(4y\pi)}{2(4\pi)^2}$  in  $\Omega = [0, 1]^2$
- We work on a sequence of Voronoi meshes with 16, 64, 100, 256, 1024, 4096 elements



Figure: The first three meshes of the sequence

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Figure: The first three meshes of the sequence

• We study the  $L^2$ ,  $H^1$  and  $L^\infty$  error with different approaches

# Convergence plots



## Limits and perspectives

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- The adaptive choice of the number of RBs in dependence on the geometry of the polygons has been done in a rough way computing the ratio between the radius of the inscribe circle with the radius of the circumscribed one
- We need to find a robust criterion to understand how many RBs we need to get a good approximation on each K and improve the convergence (Artificial Intelligence?)
- The idea is to obtain a cheap method for the post-processing of VEM solutions (for instance, reconstruction in subdomains)

## Some references

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